**University of Msila** 

FACULTY OF MATHEMATICS AND

**INFORMATICS** 

**DEPARTMENT OF COMPUTER** 

SCIENCE

# Modeling of dynamic discret events systems 2 (DDES)

#### EVENT-BASED APPROACH

In this approach, a system is modeled by defining the changes that take place when events occur. The designer must therefore determine which events can change the state of the system, and then develop the logic associated with each type of event.

#### Example:

The system to be modeled is a bank where customers arrive and are served by a single waiter (employee). Customers arrive at the bank, wait in front of a counter, are served and then leave the bank.



- he state of the system is defined by the state of the server and the number of waiting customers. In this way, the state remains constant except when a customer arrives or when a customer leaves the bank.
- The event-based approach describes what happens when a customer arrives and when a customer leaves the service. Since the system only changes state at the instants of these two events (customer arrival or customer departure), these are sufficient to fully describe the dynamics of the system.



Evénement : Fin\_De\_Service

Ejecter le client du système ;

```
Si Clients_En_Attente > 0
```

<u>Alors</u>

/\* on décrémente le nombre de clients en attente \*/ /\* on engage un nouveau client dans le service \*/

- Clients\_En\_Attente ← Clients\_En\_Attente 1
- Planifier un Evénement Fin\_De\_Service pour ce client à la date : TNOW + Durée de service du client

#### <u>Sinon</u>

/\* le serveur vient de finir de servir un client et \*/ /\* il n' y a plus de clients en attente \*/

● ETAT\_SERVEUR ← Libre

Fin

Fin événement Fin\_De\_Service





Retour au Noyau

- In the event-driven approach, the kernel's main tasks are :
   1) Time control: determine the date of the next event and initialize the clock with this date
  - 2) Identify events: determine which events must take place at the current time (clock)
    3) Execute events: trigger the events identified as due to take place.

#### **Example:**

The arrival of a customer will trigger the addition of an End\_Of\_Service event to the calendar. It also triggers the addition of a Customer\_Arrival event corresponding to the next customer.

The end of a service can trigger the addition of an End\_Of\_Service event to the schedule, if the queue is not empty, and which in this case concerns the new customer in the queue. In this way, each event in the list must include at least the following two items of information: the of the date occurrence of event identification number) the event (usually a Information on the entities involved in the event may also be useful (e.g.: selection of a customer in the queue). As the simulation the kernel will execute a 2-phase cycle: progresses,

#### 1) Time control: this phase includes

a) determining the date of the next event by examining the calendar (list of events)

b) initializing the clock with the date of the next event

c) construction of a list of current events including all events whose date of occurrence is equal to the clock.

#### 2) Execution of current events

Current events are executed under kernel control. No event can be triggered without the kernel. This ensures that the logical sequence of events is entirely controlled by the kernel. Once an event has been executed, it is deleted (from the calendar and from the list of current events).

This cycle is repeated until the end of the simulation. A simplified kernel flowchart might look like this:



#### Components of a discrete-event simulation model

The following components are found in most discrete-event simulation models adopting the "event" approach:

- System state: defined by a collection of state variables describing the state of the system at a particular time.
- Simulation clock: a variable indicating the current value of simulation time.
- Event list: A list containing the dates of occurrence of events scheduled to take place in the future.
- Statistical counters: Variables used to collect statistical information on system performance.
- Initialization routine. Subroutine used to initialize the simulation model at time zero.
- **Time management routine** (time control): subroutine which determines the next event from the event list and initializes the simulation clock with the date of occurrence of this current event.
- Event routines: subroutines that update the system state when a particular type of event occurs (there is an event routine for each type of event).
- utility routine library: set of subroutines used to generate random variables identified as part of the simulation model.
- Results generator: subroutine which calculates estimates (from statistical counters) of the desired performance measures and generates a report at the end of the simulation.
- Main program: responsible for initializing the system state at the start of the simulation, and all variables
  used during the simulation. It invokes the time management routine to determine the next event to take
  place and passes control to the routine associated with that event. It also checks whether the end of the
  simulation has been reached and, if so, invokes the results generator.



# **Simulation languages**

The first simulation languages appeared around 1960, and were more oriented towards the representation and simulation of discontinuous (discrete-event) systems.

Simulateur	Pays	Date
SIMSCRIPT	USA	1963
SIMULA	Norvège	1966
GPSS	USA	1968
GASP	USA	1974
Q-GERT	USA	1977
SLAM	USA	1979
QNAP	France	1980
CAPS/ECSL	Grande Bretagne	1980
SIMAN	USA	1982

## **Simulation languages**

**GPSS (General Purpose Simulation System)** is a language offering a process-oriented approach. It was first developed by IBM in 1961, and several versions have since been released, the most recent of which is GPSS/V. GPSS is one of the forerunners of the "process" modeling approach, and has graphical representation support like most current languages.



r r	Exemple de	Modèle en GPSS
	SIMULATE	
	GENERATE	RVEXPO(1,1.0)
	QUEUE	QSERVER
	SEIZE	SERVER
VEQ	DEPART	QSERVER
	TEST L	N\$LVEQ,1000,STOP
	ADVANCE	RVEXP0(2,0.5)
STOP	RELEASE	SERVER
r	TERMINATE	1
t	Instructions	de Controle
	START End	1000

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-Injection des entités dans le creation système entrée dans la file d'attente demande d'allocation du serveur sortie de la la file d'attente Test si l'execution est arrivée a sa fin occuper le serveur pendant t = durée de service libérer le serveur entité quitte le système

faire 1 seule simulation qui dure 1000 unités

RELATIVE CLOCK: 1014.1565	ABSOLUTE CLOCK: 1014.1565
BLOCK CURRENT	TOTAL
1	1000
2	1000
3	1000
LVEQ	1000
5	1000
6	999
STOP	1000
8	1000

#### --AVG-UTIL-DURING --

FACILITY	TOTAL	AVAIL	UNAVL	ENTRIES	AVERAGE	CURRENT	PERCENT	SEIZING	PREEMPTING
	TIME	TIME	TIME		TIME/XACT	STATUS	AVAIL	XACT	XACT
SERVER	0.516			1000	0.523	AVAIL			

QUEUE	MAXIMUM	AVERAGE	TOTAL	ZERO	PERCENT	AVERAGE	SAVERAGE	QTABLE	CURRENT
	Contents	CONTENTS	Entries	Entries	ZEROS	TIME/UNIT	TIME/UNIT	NUMBER	CONTENTS
SERVERQ	8	0.605	1000	454	45.4	0.614	1.124		0

RAMDOM	ANTITHETIC	INITIAL	CURRENT	SAMPLE	CHI-SQUARE
STREAM	VARIATES	Position	POSITION	COUNT	UNIFORMITY
1	OFF	100000 200000	101001 200999	1001 999	0.71 0.69

### **Specialized simulators**

A specialized simulator is a simulation tool offering a description language (or an interactive data input system) in which the instructions (or primitives) are objects of the system to be modeled. The primitives offered are, in this case, aggregates of elementary processes (queues, resources, activities, etc ...) representing a particular object of the system (machine, stock, conveyor, pallet, etc ... in the case of a production system) and its behavior.
 One of the advantages of specialized simulators is the reduced conceptualization effort required in the modeling stage. The following figure shows how easy it is to represent the same problem (here, the modeling of a diverging switch in a handling network) using a specialized "handling network" simulator, and a general simulation language (here, SLAM II).



### **Specialized simulators**

Simulateur	Sociétés	Domaine
MAP/1	Pritsker Associates - USA	systèmes de production
SIMFLEX	INRIA/SIMULOG - France	Ateliers Flexibles
SAME/AGVS	SERI/RENAULT Automation	Circuits de manutention
SIMUFLEX	RAMSES Automation - France	Ateliers Flexibles
NETWORK II.5	CACI-Los Angeles - USA	Réseaux locaux LAN
COMNET II.5	CACI-Los Angeles - USA	Réseaux Télécom WAN