

Binary Relations on a Set

3.1 Basic Definitions

Definition 3.1 (Binary Relation) Let E be a set. A binary relation \mathcal{R} on E is a property that applies to pairs of elements from E . We denote $x\mathcal{R}y$ to indicate that the property is true for the pair $(x, y) \in E \times E$.

Example

1. The inequality \leq is a relation on \mathbb{N} , \mathbb{Z} , and \mathbb{R} .
2. The inclusion relation in the power set of E : $A\mathcal{R}B \Leftrightarrow A \subset B$.
3. The divisibility relation on the integers: $m\mathcal{R}n \Leftrightarrow m$ divides n .

Definition 3.2 Let \mathcal{R} be a relation on a set E .

1. \mathcal{R} is reflexive if for every $x \in E$, $x\mathcal{R}x$ holds.
2. \mathcal{R} is symmetric if for all $x, y \in E$, $x\mathcal{R}y \Rightarrow y\mathcal{R}x$.
3. \mathcal{R} is antisymmetric if for all $x, y \in E$, $(x\mathcal{R}y \wedge y\mathcal{R}x) \Rightarrow x = y$.
4. \mathcal{R} is transitive if for all $x, y, z \in E$, $(x\mathcal{R}y \wedge y\mathcal{R}z) \Rightarrow x\mathcal{R}z$.

3.2 Equivalence Relations

Definition 3.3 (Equivalence Relation) A binary relation \mathcal{R} on E is an equivalence relation if and only if it is reflexive, symmetric, and transitive.

Example 1 The relation \mathcal{R} of "being parallel" is an equivalence relation for the set E of affine lines in the plane:

1. Reflexivity: A line is parallel to itself.
2. Symmetry: If line D is parallel to D' , then D' is parallel to D .
3. Transitivity: If line D is parallel to D' and D' is parallel to D'' , then D is parallel to D'' .

Example 2 Consider the following relation on \mathbb{Z} :

$$x\mathcal{R}y \Leftrightarrow \exists k \in \mathbb{Z} \mid x - y = 2k$$

1. \mathcal{R} is reflexive because $\exists k = 0 \mid x - x = 2k = 0$, thus $x\mathcal{R}x$.
2. Suppose $x\mathcal{R}y$, then $\exists k \in \mathbb{Z} \mid x - y = 2k \Rightarrow y - x = 2k'$ with $k' = -k \in \mathbb{Z}$. Therefore, $y\mathcal{R}x$. Hence, \mathcal{R} is symmetric.
3. Suppose $x\mathcal{R}y$ and $y\mathcal{R}z$. Then, $(\exists k \in \mathbb{Z} \mid x - y = 2k)$ and $(\exists k' \in \mathbb{Z} \mid y - z = 2k')$ by adding these equations, we obtain $x - z = 2k''$ with $k'' = (k + k') \in \mathbb{Z}$. Thus, $x\mathcal{R}z$. Therefore, \mathcal{R} is transitive. Consequently, \mathcal{R} is an equivalence relation.

Definition 3.4 Let \mathcal{R} be an equivalence relation on a set E . The equivalence class of an element $x \in E$ is the set of elements in E that are related to x by \mathcal{R} , denoted by $\mathcal{C}(x)$ or \bar{x} :

$$\bar{x} = \{y \in E \mid y\mathcal{R}x\}$$

Definition 3.5 Let \mathcal{R} be an equivalence relation on a set E . The quotient set of E by \mathcal{R} is the set of equivalence classes of \mathcal{R} , denoted by E/\mathcal{R} :

$$E/\mathcal{R} = \{\bar{x} \mid x \in E\}$$

Example In the previous example, we have

$$\begin{aligned}\bar{x} &= \{y \in E \mid y\mathcal{R}x\} \\ &= \{y \in E \mid x - y = 2k\} \\ &= \{x - 2k : k \in \mathbb{Z}\} \\ &= \{\dots, x - 4, x - 2, x, x + 2, x + 4, \dots\}.\end{aligned}$$

$$\bar{0} = \{y \in E \mid 0\mathcal{R}y\} = \{\dots, -4, -2, 0, 2, 4, \dots\}, \bar{1} = \{y \in E \mid 1\mathcal{R}y\} = \{\dots, -3, -1, 1, 3, \dots\}$$

and $\bar{2} = \bar{0}$. Therefore, $\mathbb{Z}/\mathcal{R} = \{\bar{x} \mid x \in E\} = \{\bar{0}, \bar{1}\}$

Proposition 3.1 Let \mathcal{R} be an equivalence relation on E . Then

1. An equivalence class is a subset of the set E , i.e., for all $x \in E$, $\bar{x} \subset E$.
2. An equivalence class is never empty, i.e., for all $x \in E$, $\bar{x} \neq \phi$.
3. The intersection of two distinct equivalence classes is empty, i.e., for all $x, y \in E$, $\bar{x} \cap \bar{y} = \phi$.
4. For all $x, y \in E$, $x\mathcal{R}y \Leftrightarrow \bar{x} = \bar{y}$.

Theorem 3.1 Let \mathcal{R} be an equivalence relation on E . The equivalence classes $(\bar{x})_{x \in E}$ form a partition of E :

$$E = \cup_{x \in E} \bar{x}$$

3.3 Order Relation

Definition 3.6 (Order Relation) A binary relation \mathcal{R} on E is an order relation if and only if it is reflexive, antisymmetric, and transitive. We then say that (E, \mathcal{R}) is an ordered set.

Example.

1. The inequality \leq is an order relation on \mathbb{N} , \mathbb{Z} , and \mathbb{R} .
2. The inclusion relation in the power set of E is an order relation: $A\mathcal{R}B \Leftrightarrow A \subset B$.

Definition 3.7 Let \mathcal{R} be an order relation on E . Two elements x and y of E are said to be comparable if $x\mathcal{R}y$ or $y\mathcal{R}x$.

Definition 3.8 (Total Order and Partial Order) Let \mathcal{R} be an order relation on E . If any two elements x and y are always comparable, we say that \mathcal{R} is a total order relation and the set E is called totally ordered. Otherwise (i.e., if there exist at least two non-comparable elements x and y), we say that \mathcal{R} is a partial order relation and the set E is called partially ordered.

Example.

1. \leq is a total order on \mathbb{N} , \mathbb{Z} , and \mathbb{R} .
2. The divisibility relation in \mathbb{N}^* is a partial order.

Definition 3.9 Let \mathcal{R} be an order relation on E , and let M, m be two elements of E .

1. M is an upper bound of a subset A of E if $x\mathcal{R}M$ for every $x \in A$.
2. m is a lower bound of a subset A of E if $m\mathcal{R}x$ for every $x \in A$.

Example.

1. The set $\{8, 10, 12\}$ is bounded above by 120 and bounded below by 2 for the divisibility relation $"/$ on \mathbb{N} .
2. $\mathcal{P}(E)$ is bounded below by \emptyset and bounded above by E for the inclusion relation \subset .

3.4 Exercises with Solutions

Exercise 1. In \mathbb{R} , the binary relation \mathcal{R} is defined as follows:

$$\forall x, y \in \mathbb{R} : x\mathcal{R}y \iff x^2 - 1 = y^2 - 1$$

1. Show that \mathcal{R} is an equivalence relation on \mathbb{R} .
2. Determine the quotient set \mathbb{R}/\mathcal{R} .

Exercise 2. For every $n \in \mathbb{N}^*$, a binary relation on \mathbb{Z} is defined by

$$\forall x, y \in \mathbb{Z} : x\mathcal{R}y \iff \exists k \in \mathbb{Z} \mid x - y = kn$$

1. Show that \mathcal{R} is an equivalence relation on \mathbb{Z} .
2. Assume that $n = 3$:
 - (a) Determine the equivalence class of $x \in \mathbb{Z}$. Deduce the classes $\bar{0}, \bar{1}, \bar{2}$.
 - (b) Show that $\forall m \in \mathbb{Z} : \bar{0} = \overline{3m}, \bar{1} = \overline{3m+1}, \bar{2} = \overline{3m+2}$.
 - (c) Show that $\bar{0} \cap \bar{1} = \emptyset, \bar{1} \cap \bar{2} = \emptyset, \bar{0} \cap \bar{2} = \emptyset$. Deduce the quotient set \mathbb{Z}/\mathcal{R} .

Exercise 3. Let E be a set and let A be a subset of E . A binary relation \mathcal{R} is defined on $\mathcal{P}(E)$ as follows:

$$\forall X, Y \in \mathcal{P}(E) : X\mathcal{R}Y \iff A \cap X = A \cap Y$$

1. Show that \mathcal{R} is an equivalence relation on $\mathcal{P}(E)$.
2. Determine the equivalence classes of \emptyset and E . Deduce \bar{A} and $\overline{C_E(A)}$.

Exercise 4. Let \mathcal{R} be a binary relation on \mathbb{R}^3 defined by

$$(x, y, z)\mathcal{R}(a, b, c) \iff (|x - a| \leq b - y \text{ and } z = c).$$

1. Show that \mathcal{R} is a partial order relation on \mathbb{R}^3 .
2. Is the order total on \mathbb{R}^3 ?

Exercise 5. A binary relation \mathcal{R} is defined on \mathbb{R}^2 as follows:

$$\forall (x_1, y_1), (x_2, y_2) \in \mathbb{R}^2 : (x_1, y_1)\mathcal{R}(x_2, y_2) \iff x_1 \leq x_2 \text{ and } y_1 \leq y_2.$$

1. Show that \mathcal{R} is an order relation on \mathbb{R}^2 .
2. Are the elements $(2, 4), (3, 1)$ of \mathbb{R}^2 comparable by \mathcal{R} ?
3. Is the order total on \mathbb{R}^2 ?
4. Determine the set of upper bounds of $A = \{(1, 2), (3, 1)\} \subset \mathbb{R}^2$.

Exercise 6. Determine whether the following relations \mathcal{R} are order relations:

1. $\forall x, y \in \mathbb{R} : x\mathcal{R}y \iff e^x \leq e^y$;
2. $\forall x, y \in \mathbb{R} : x\mathcal{R}y \iff |x| \leq |y|$;
3. $\forall x, y \in \mathbb{N} : x\mathcal{R}y \iff \exists p, q \geq 1 \mid y = px^q$ (where p and q are integers);
4. $\forall x, y \in \mathbb{N}^* : x\mathcal{R}y \iff \exists m \in \mathbb{N}^* \mid y = mx$;
5. $\forall x, y \in]1, +\infty[: x\mathcal{R}y \iff \frac{x}{1+x^2} \geq \frac{y}{1+y^2}$.

3.4.1 Solution

Exercise 1.

$$1. \forall x, y \in \mathbb{R} : x\mathcal{R}y \iff x^2 - 1 = y^2 - 1$$

(i) Reflexivity: $\forall x \in \mathbb{R}, x^2 - 1 = x^2 - 1 \Rightarrow xRx$.

(ii) Symmetry: $x\mathcal{R}y \iff x^2 - 1 = y^2 - 1 \Rightarrow y^2 - 1 = x^2 - 1 \Rightarrow yRx$.

(iii) Transitivity:

$$\begin{cases} x\mathcal{R}y \\ y\mathcal{R}z \end{cases} \iff \begin{cases} x^2 - 1 = y^2 - 1 \\ y^2 - 1 = z^2 - 1 \end{cases} \Rightarrow x^2 - 1 = z^2 - 1 \Rightarrow x\mathcal{R}z.$$

Therefore, \mathcal{R} is an equivalence relation.

$$2. \mathbb{R}/\mathbb{R} = \{\bar{x} : x \in \mathbb{R}\}.$$

We have $\bar{x} = \{y \in \mathbb{R} \mid yRx\} = \{y \in \mathbb{R} \mid y^2 - 1 = x^2 - 1\} = \{x, -x \mid x \in \mathbb{R}\}$

Thus, $\mathbb{R}/\mathbb{R} = \{\{x_1 - x\}, x \in \mathbb{R}\}$.

Exercise 2.

$$1. \forall x, y \in \mathbb{Z} : x\mathcal{R}y \iff \exists k \in \mathbb{Z} \mid x - y = kn.$$

- Reflexivity: We know that $\forall x \in \mathbb{Z} : x - x = 0 = 0 \cdot n$ with $k = 0 \in \mathbb{Z}$, so xRx .

- Symmetry: $x\mathcal{R}y \iff x - y = kn \Rightarrow y - x = (-k) \cdot n = k' \cdot n$ with $k' = -k \in \mathbb{Z}$. Thus,

yRx .

- Transitivity:

$$\begin{cases} x\mathcal{R}y \\ y\mathcal{R}z \end{cases} \Leftrightarrow \begin{cases} x - y = k_1 \cdot n/k_1 \in \mathbb{Z} \\ y - z = k_2 \cdot n/k_2 \in \mathbb{Z} \end{cases} \quad ; \text{ Summing both sides:}$$

$$x - z = (k_1 + k_2)n = k_3 \cdot n \text{ with } k_3 = k_1 + k_2 \in \mathbb{Z}$$

Therefore, $x\mathcal{R}z$

2. For $n = 3 : \forall x, y \in \mathbb{Z} : x\mathcal{R}y \Leftrightarrow \exists k \in \mathbb{Z} : x - y = 3k$.

(a) For any

$$\begin{aligned} x \in \mathbb{Z} : \bar{x} &= \{y \in \mathbb{Z} : y\mathcal{R}x\} = \{y \in \mathbb{Z} : y = x + 3k\} \\ &= \{x + 3k \mid k \in \mathbb{Z}\}. \end{aligned}$$

In particular:

$$\bar{0} = \{y \in \mathbb{Z} : y\mathcal{R}0\} = \{3k \mid k \in \mathbb{Z}\} = 3\mathbb{Z}$$

$$\bar{1} = \{y \in \mathbb{Z} : y\mathcal{R}1\} = \{3k + 1 \mid k \in \mathbb{Z}\} = 3\mathbb{Z} + 1$$

$$\bar{2} = \{y \in \mathbb{Z} : y\mathcal{R}2\} = \{3k + 2 \mid k \in \mathbb{Z}\} = 3\mathbb{Z} + 2.$$

(b)

For all $m \in \mathbb{Z}$:

$$\begin{cases} \bar{0} = 3\bar{m} \\ \bar{1} = 3\bar{m} + 1 \\ \bar{2} = 3\bar{m} + 2 \end{cases} \quad \text{because } \forall m \in \mathbb{Z} : \begin{cases} 0\mathcal{R}(3m) \\ 1\mathcal{R}(3m + 1) \\ 2\mathcal{R}(3m + 2) \end{cases} .$$

Indeed, for all $m \in \mathbb{Z}$:

$$\begin{cases} 0 - (3m) = 3(-m) \\ 1 - (3m + 1) = 3(-m) \\ 2 - (3m + 2) = 3(-m) \end{cases}, \quad -m \in \mathbb{Z} .$$

(C)

We have:

$$\begin{cases} \bar{0} \cap \bar{1} = \emptyset \\ \bar{1} \cap \bar{2} = \emptyset \\ \bar{0} \cap \bar{2} = \emptyset \end{cases}, \text{ because } \begin{cases} 0 \not\mathcal{R} 1 \\ 1 \not\mathcal{R} 2 \\ 0 \not\mathcal{R} 2 \end{cases}. \text{ Indeed, } \begin{cases} 0 - 1 = -1 \neq 3k_1 \\ 1 - 2 = -1 \neq 3k_2 \\ 0 - 2 = -2 \neq 3k_3 \end{cases}, \quad k_1, k_2, k_3 \in \mathbb{Z}.$$

We know that:

$$\begin{aligned} \mathbb{Z}/R &= \{\bar{x} : x \in \mathbb{Z}\} \\ &= \{\bar{x} : x = 3m\} \cup \{\bar{x} : x = 3m + 1\} \cup \{\bar{x} : x = 3m + 2\}. \\ &= \{\bar{0}, \bar{1}, \bar{2}\}. \end{aligned}$$

Exercise 4. $(x, y, z)R(a, b, c) \Leftrightarrow (|x - a| \leq b - y \text{ and } z = c)$

(1)

(i) Reflexivity: $(x, y, z)R(x, y, z) \Leftrightarrow (|x - x| = 0 \leq y - y = 0 \text{ and } z = z)$, hence R is reflexive.

(ii) Anti-symmetry: Suppose $(v, y, z)R(a, b, c)$ and $(a, b, c)R(x, y, z)$

This implies $[(|x - a| \leq b - y \quad (*) \text{ and } |a - x| \leq y - b \quad (**)) \text{ and } z = c]$

Then, $(*) + (**)$ gives: $x = a$, replacing $x = a$ in $(*)$ and $(**)$ we find $y = b$. Thus, $(x, y, z) = (a, b, c)$. Therefore, R is anti-symmetric.

(iii) Transitivity: Suppose $(v, y, z)R(a, b, c)$ and $(a, b, c)R(\alpha, \beta, \gamma)$

This implies $[(|x - a| \leq b - y \quad (*) \text{ and } |a - \alpha| \leq \beta - b \quad (**)) \text{ and } z = c = \gamma]$

Thus, $(*) + (**)$ gives $(|x - a| + |a - \alpha| \leq b - y + \beta - b \text{ and } z = c = \gamma)$.

And since $(|x - \alpha| = |x - a + a - \alpha| \leq |x - a| + |a - \alpha| \leq y + \beta \text{ and } z = \gamma)$ implies $(x, y, z)R(\alpha, \beta, \gamma)$. Hence, R is transitive.

Therefore, R is a partial order relation on \mathbb{R}^3 .

(2) R is not total because $\exists(x, y, z) = (0, 0, 2) \in \mathbb{R}^3$ and $(a, b, c) = (0, 0, 3) \in \mathbb{R}^3$ such that $(0, 0, 2) \not\mathcal{R} (0, 0, 3)$ and $(0, 0, 3) \not\mathcal{R} (0, 0, 2)$.

Exercise 5. $\forall (x_1, y_1), (x_2, y_2) \in \mathbb{R}^2 : x_1 \leq x_2 \text{ and } y_1 \leq y_2$.

(1)

(i) Reflexivity: We know that

$$\forall (x, y) \in \mathbb{R}^2 : \begin{cases} x \leq x \\ y \leq y \end{cases} \Rightarrow (x, y)R(x, y) \Rightarrow R \text{ is reflexive.}$$

(ii) Anti-symmetry: Suppose $(x_1, y_1)R(x_2, y_2)$ and $(x_2, y_2)R(x_1, y_1)$

$$\Rightarrow \begin{cases} x_1 \leq x_2 \wedge y_1 \leq y_2 \\ \quad \quad \quad \wedge \\ x_2 \leq x_1 \wedge y_2 \leq y_1 \end{cases} \Rightarrow \begin{cases} x_1 = x_2 \\ \quad \quad \quad \wedge \\ y_1 = y_2 \end{cases} \Rightarrow (x_1, y_1) = (x_2, y_2). \text{ Thus, } R \text{ is anti-symmetric.}$$

(iii) Transitivity: Let $(x_1, y_1), (x_2, y_2), (x_3, y_3) \in \mathbb{R}^2$

$$\begin{cases} (x_1, y_1)R(x_2, y_2) \\ \quad \quad \quad \wedge \\ (x_2, y_2)R(x_3, y_3) \end{cases} \Rightarrow \begin{cases} x_1 \leq x_2 \wedge y_1 \leq y_2 \\ \quad \quad \quad \wedge \\ x_2 \leq x_3 \wedge y_2 \leq y_3 \end{cases} \Rightarrow \begin{cases} x_1 \leq x_3 \\ \quad \quad \quad \wedge \\ y_1 \leq y_3 \end{cases} \Rightarrow (x_1, y_1)R(x_3, y_3)$$

Therefore, R is transitive. Hence, R is a partial order relation on \mathbb{R}^2 .(2) $(2, 4)$ and $(3, 1)$ are not comparable because $(1, 4)$ and $(3, 1)$ do not satisfy the relation. In

$$\text{fact, } \begin{cases} 2 \leq 3 \\ \quad \quad \quad \wedge \\ 4 \not\leq 1 \end{cases} \text{ and } \begin{cases} 3 \not\leq 2 \\ \quad \quad \quad \wedge \\ 1 \leq 2 \end{cases} \Rightarrow \begin{cases} (2, 4) \not R (3, 1) \\ \quad \quad \quad \wedge \\ (3, 1) \not R (2, 4) \end{cases}$$

(3) The order is partial because $\exists a = (2, 4)$ and $b = (3, 1)$ where $a \not R b$ and $b \not R a$.(4) $t = (x, y) \in \mathbb{R}^2$ is an upper bound of A if $\forall a \in A : aRt$.

$$\Rightarrow \begin{cases} (1, 2)R(x, y) \\ \quad \quad \quad \wedge \\ (3, 1)R(x, y) \end{cases} \Rightarrow \begin{cases} 1 \leq x \wedge 2 \leq y. \\ \quad \quad \quad \wedge \\ 3 \leq x \wedge 1 \leq y. \end{cases} \Rightarrow \begin{cases} x \geq 3 \\ \quad \quad \quad \wedge \\ y \geq 2 \end{cases}$$

$$\Rightarrow \text{Maj}_{\mathbb{R}^2}(A) = \{(x, y) : x \geq 3 \wedge y \geq 2\}.$$

Algebraic Structures

4.1 Internal Composition Laws and Their Properties

4.1.1 Internal Composition Laws

Definition 4.1 Let E be a set. An internal composition law $*$ on E is a mapping from $E \times E$ to E :

$$\begin{aligned} * : E \times E &\longrightarrow E \\ (x, y) &\mapsto x * y \end{aligned}$$

Notations

1. Instead of "internal composition law," we also say "operation of internal composition" or simply "internal operation."
2. $(E, *)$ is often used to denote a set E equipped with an internal operation $*$.

Example.

1. The laws \cup (union), \cap (intersection), and Δ (symmetric difference) on $\mathcal{P}(E)$ (the power set of E).
2. The law (composition) on $\mathcal{F}(E)$ (the set of functions from E to E).

3. The laws $+$ and \times on $\mathbb{N}, \mathbb{Z}, \mathbb{Q}, \mathbb{R}$, and \mathbb{C} .
4. Let $*$ be defined on \mathbb{R} by $x * y = \frac{1}{x+y}$. Then $*$ is not an internal operation since $(-1, 1)$ does not have an image.

Definition 4.2 (Stable Subset for an Operation) Let E be a set equipped with an internal composition law $*$ and F be a subset of E . We say that F is stable under the law $*$ if

$$\forall (x, y) \in F \times F : x * y \in F$$

Example.

1. \mathbb{R}^+ and \mathbb{R}^- are two stable subsets of \mathbb{R} under the operation $+$.
2. For the operation \times , \mathbb{R}^+ is still a stable subset, but \mathbb{R}^- is not.

4.1.2 Properties of internal composition laws

Definition 4.3 (Commutativity and Associativity) Let E be a set equipped with an internal composition law $*$.

We say that $*$ is commutative if $\forall (x, y) \in E^2 : x * y = y * x$.

We say that $*$ is associative if $\forall (x, y, z) \in E^3 : (x * y) * z = x * (y * z)$.

Example.

1. The addition and multiplication laws on $\mathbb{N}, \mathbb{Z}, \mathbb{Q}, \mathbb{R}$, and \mathbb{C} are commutative and associative.
2. Also, the union (\cup), intersection (\cap), and symmetric difference (Δ) laws on $\mathcal{P}(E)$ are commutative and associative.
3. The composition law (\circ) on $\mathcal{F}(E)$ is associative but not commutative, because $f \circ g \neq g \circ f$ in general.
4. Let $*$ be the law defined on \mathbb{Q} by: $x * y = \frac{x+y}{2}$. Then $*$ is commutative, because $x * y = \frac{x+y}{2} = \frac{y+x}{2} = y * x$, but it is not associative,

because $(-1 * 0) * 1 = \frac{1}{4} \neq -1 * (0 * 1) = \frac{-1}{4}$.

Definition 4.4 (Neutral Element) Let E be a set equipped with an internal composition law $*$. Let e be an element of E . We say that e is the neutral element for the law $*$, if

$$\forall x \in E : x * e = e * x = x$$

Remark 4.1 If the law $*$ is commutative, the equality $x * e = e * x$ is automatically satisfied.

Example.

1. In $\mathbb{N}, \mathbb{Z}, \mathbb{Q}, \mathbb{R}$, and \mathbb{C} , 0 is neutral for the addition law, and 1 is neutral for the multiplication law.
2. In $\mathcal{P}(E)$, the empty set (\emptyset) is neutral for the union law (\cup), and E is neutral for the intersection law (\cap).
3. Let $*$ be the law defined on \mathbb{R} by: $x * y = x + y - 1$. Then $e = 1$ is a neutral element, because $x * e = x \Rightarrow x + e - 1 = x$. Thus, $e = 1$.

Proposition 4.1 (Uniqueness of the Neutral Element) The neutral element of E for the law $*$ if it exists, is unique.

Proof. Indeed, let e' be another neutral element for $*$, then $e' = e' * e = e * e' = e$. Thus, the neutral element is unique.

Definition 4.5 (Inverse Element) Let E be a set equipped with an internal composition law $*$ and let e be a neutral element. We say that the element x of E has an inverse element x' of E , if $\forall x \in E : x * x' = x' * x = e$.

Example.

1. In \mathbb{R} , the invertible elements for the multiplication law (\times) are the non-zero elements.
2. Let $*$ be the law defined on \mathbb{R} by: $x * y = x + y - 1$. Then $x \in \mathbb{R}$ has an inverse element $x' = 2 - x$, because $x * x' = 1 \Rightarrow x + x' - 1 = 1$. Thus, $x' = 2 - x$.

4.1.3 Properties of internal composition laws

Definition 4.3 (Commutativity and Associativity) Let E be a set equipped with an internal composition law $*$.

We say that $*$ is commutative if $\forall(x, y) \in E^2 : x * y = y * x$.

We say that $*$ is associative if $\forall(x, y, z) \in E^3 : (x * y) * z = x * (y * z)$.

Example.

1. The addition and multiplication laws on \mathbb{N} , \mathbb{Z} , \mathbb{Q} , \mathbb{R} , and \mathbb{C} are commutative and associative.
2. The union (\cup), intersection (\cap), and symmetric difference (Δ) laws on $\mathcal{P}(E)$ are commutative and associative.
3. The composition law (\circ) on $\mathcal{F}(E)$ is associative but not commutative, because $f \circ g \neq g \circ f$ in general.
4. Let $*$ be the law defined on \mathbb{Q} by: $x * y = \frac{x+y}{2}$. Then $*$ is commutative, because $x * y = \frac{x+y}{2} = \frac{y+x}{2} = y * x$, but it is not associative, because $(-1 * 0) * 1 = \frac{1}{4} \neq -1 * (0 * 1) = \frac{-1}{4}$.

Definition 4.4 (Neutral Element) Let E be a set equipped with an internal composition law $*$. Let e be an element of E . We say that e is the neutral element for the law $*$ if $\forall x \in E : x * e = e * x = x$.

Remark 4.1 If the law $*$ is commutative, the equality $x * e = e * x$ is automatically satisfied.

Example.

1. In \mathbb{N} , \mathbb{Z} , \mathbb{Q} , \mathbb{R} , and \mathbb{C} , 0 is the neutral element for the addition law, and 1 is the neutral element for the multiplication law.
2. In $\mathcal{P}(E)$, the empty set \emptyset is the neutral element for the union law \cup , and E is the neutral element for the intersection law \cap .