

Final Exam (Solution with Scale)

(Scale: 20 Points, Duration: 1 Hour and 30 Minutes)

Single Choice Questions: (4 pts)

Check the *only* correct answer.

1. Binary Search is not considered as a pure	2. Which of the following is a FALSE statement
Divide and Conquer method because:	about a spanning tree of a graph G?
A. It has no worst-case scenario	\Box A. It is a subgraph of G
B. Only one sub-array is conquered after each	B. It includes every vertex of G
division (0.5)	🔀 C. It can be cyclic (0.5)
C. It is not efficient for large input sizes	D. None of the above
\Box D. All of the above	
3. What is the main drawback of an ideal	4. Which of the following is an adaptive sorting
hashing?	algorithm?
A. Causes collision problem	🔀 A. Bubble Sort (0.5)
B. Requires too much space (0.5)	□ B. Selection Sort
C. Both A and B	🗆 C. Merge Sort
D. None of the above	\Box D. All of the above
5. Which problem from the following can be	6. If we want to analyze the elements frequency
easily solved if the elements of a list are sorted?	of a list, which technique would make the
A. Searching for an element	problem easier?
B. Identifying largest/smallest element	🛛 A. Sorting the list (0.5)
□ C. Detecting duplicate values	B. Divide and Conquer
D. All of the above (0.5)	C. Recursion
	D. None of the above
7. The time complexity of a sorting algorithms	8. Is it possible to make Selection Sort adaptive?
depends mostly on:	🗆 A. Yes
🗆 A. Adaptability	🔀 B. No (0.5)
□ B. Number of swaps	
C. Number of comparisons (0.5)	
□ D. All of the above	

Exercise 1: (6 pts)

1. For arr = [3, 3, 5, 6, 6, 7, 7, 8, 11, 12] and target = 21, algo it returns False: (1.0)

Index	0	1	2	3	4	5	6	7	8	9	
Elements	3	3	5	6	6	7	7	8	11	12	
Iteration #1	left									right	s = 15
Iteration #2		left								right	s = 15
Iteration #3			left							right	s = 17
Iteration #4				left						right	s = 18
Iteration #5					left					right	s = 18
Iteration #6						left				right	s = 19
Iteration #7							left			right	s = 19
Iteration #8								left		right	s = 20
Iteration #9									left	right	s = 23
End.									left		
									=		
									right		

For arr = [3, 3, 5, 6, 6, 7, 7, 8, 11, 12] and target = 13, algo returns **True**. (1.0)

Index	0	1	2	3	4	5	6	7	8	9	
Elements	3	3	5	6	6	7	7	8	11	12	
Iteration #1	left									right	s = 15
Iteration #2	left								right		s = 14
Iteration #3	left							right			s = 11
Iteration #4		left						right			s = 11
Iteration #5			left					right			s = 13
End.											

algo checks if there is any pair in the array that sums up to the given target. (0.5)

2. Total number of primitive instructions: $T(n) = \frac{13}{2}n + \frac{5}{2} = O(n)$ (1.5)

At each iteration, either left index is incremented or right index is decremented. So, the loop will be executed at most n - 1 times.

3. A naïve algorithm: $T(n) = O(n^2)$ (1.5)

```
def algo2(arr, target):
for i in range(len(arr)-1):
    for j in range(i+1, len(arr)):
        if arr[i]+arr[j] == target:
            return True
return False
```

4. algo is more time efficient since $O(n) < O(n^2)$ (0.5)

Exercise 2: (5 pts)

1. Resultant HT: (1.0)

i	0	1	2	3	4	5	6	7	8	9	10	11	12
HT(i)	26	11			30							37	24

- 2. The hash table size is not sufficient (1.0). In this case, there are two solutions:
 - We increase the size of the hash table. In this case, all the elements must be inserted again. It is common to double the size of the hash table. (0.5)
 - We use separate chaining. (0.5)

3. Resultant new HT: (2.0)

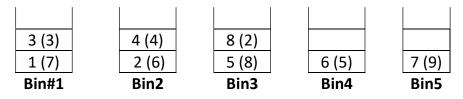
Exercise 3: (5 pts)

The problem is better known as Bin Packing Problem.

1. Strategy 1 is sequential; Strategy 2 is greedy. (1.0 pt)

2.

Following strategy 1, 5 boxes are required. (0.75 pt)



Following the strategy 2, 4 boxes are required, only. (0.75 pt)

8 (2)	3 (3)	4 (4)	6 (5)
7 (9)	5 (8)	1 (7)	2 (6)
Bin1	Bin2	Bin3	Bin4

3. The second strategy gives an optimal solution as only 4 boxes are required. (1.0 pt)

4. In the second case, the strategy 2 gives a solution with 3 boxes required: (0.75 pt)

	5 (3)	
1 (5)	3 (3)	
2 (6)	4 (5)	6 (2)
Bin1	Bin2	Bin3

However, this is not an optimal solution since there is another one with 2 boxes required only: (0.75 pt)

5 (3)	6 (2)
3 (3)	4 (5)
2 (6)	1 (5)
Bin1	Bin2